MM Optimization Algorithms

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LECTURE 3: KEY INEQUALITIES FOR MM (PART II)

QUADRATIC UPPER BOUND PRINCIPLE

Quadratic Upper Bound Principle

key mechanisms

- majorization via gradient Lipschitz continuity
- majorization via bounded Hessian
- minorization via strong convexity

Gradient Lipschitz Continuity



► f is differentiable

gradient Lipschitz continuous with constant L, i.e.,

$$\|\nabla f(x) - \nabla f(y)\| \le L||x - y|| \text{ for all } x, y \tag{1}$$

▶ from (44) ¹

$$f(x) = f(y) + \int_{0}^{1} \nabla f(y + t(x - y))^{\mathsf{T}}(x - y)dt$$
(2)
$$= f(y) + \nabla f(y)^{\mathsf{T}}(x - y) + \int_{0}^{1} \left[\nabla f(y + t(x - y)) - \nabla f(y) \right]^{\mathsf{T}}(x - y)dt$$
(3)
$$\leq f(y) + \nabla f(y)^{\mathsf{T}}(x - y) + \int_{0}^{1} \|\nabla f(y + t(x - y)) - \nabla f(y)\| \|(x - y)\| dt$$
(4)
$$\leq f(y) + \nabla f(y)^{\mathsf{T}}(x - y) + L \|x - y\|^{2} \int_{0}^{1} t dt$$
(5)
$$= f(y) + \nabla f(y)^{\mathsf{T}}(x - y) + (L/2) \|x - y\|^{2} = g(x|y)$$
(6)

¹See page 34 and substitute a = y and b = x - y.

Bounded Hessian



► *f* is twice differentiable

▶ *f* has bounded Hessians, i.e.,

$$\exists B \succ 0 \text{ s.t } B - \nabla^2 f(x) \succeq 0 \text{ for all } x \tag{7}$$

From (46)²

$$f(x) = f(y) + \nabla f(y)^{\mathsf{T}}(x-y) + \qquad (8)$$

$$\int_{0}^{1} \int_{0}^{t} (x-y)^{\mathsf{T}} \nabla^{2} f(y + \tau(x-y))(x-y) d\tau dt \qquad (9)$$

$$\leq f(y) + \nabla f(y)^{\mathsf{T}}(x-y) +$$

$$\int_{0}^{1} \int_{0}^{t} (x-y)^{\mathsf{T}} B(x-y) d\tau dt \qquad (10)$$

$$= f(y) + \nabla f(y)^{\mathsf{T}}(x-y) + (1/2)(x-y)^{\mathsf{T}} B(x-y) \qquad (11)$$

$$= g(x|y) \qquad (12)$$

²See page 35 and substitute a = y and b = x - y.

Strong Convexity

suppose

f is twice differentiable and strongly convex

thus,

$$\exists m > 0 \text{ s.t } \nabla^2 f(x) - mI \succeq 0 \text{ for all } x$$
 (13)

From (46)

$$f(x) = f(y) + \nabla f(y)^{\mathsf{T}}(x - y) + \qquad (14)$$

$$\int_{0}^{1} \int_{0}^{t} (x - y)^{\mathsf{T}} \nabla^{2} f(y + \tau(x - y))(x - y) d\tau dt \quad (15)$$

$$\geq f(y) + \nabla f(y)^{\mathsf{T}}(x - y) + \qquad m \int_{0}^{1} \int_{0}^{t} (x - y)^{\mathsf{T}}(x - y) d\tau dt \quad (16)$$

$$= f(y) + \nabla f(y)^{\mathsf{T}}(x - y) + (m/2) ||x - y||^{2} \quad (17)$$

$$= g(x|y) \quad (18)$$

EXAMPLES

Landweber's Method

- consider a positive definite matrix $A \in \mathbb{S}^n$
- we have to find the solution x^* for Ax = -b
- the task requires order of n^3 flops
 - A^{-1} is to be computed
 - \blacktriangleright if n is very large the task may be computationally challending

 \blacktriangleright but we may rely on MM principle \rightarrow avoid matrix inversion



$$x^* = \operatorname*{arg\,min}_x f(x) = (1/2) \ x^\mathsf{T} A x + b^\mathsf{T} x$$
 (19)

▶ moreover, we have ³

$$\|\nabla f(x) - \nabla f(y)\| \le \lambda_{\max}(A) \|x - y\|$$
(20)
$$\le \|A\|_* \|x - y\|$$
(21)

where $\|\cdot\|_*$ is any matrix norm

³See p. 497, *Matrix Analysis and Applied Linear Algebra* by C. D. Meyer, 2000.

▶ thus, with any $\rho > ||A||_*$, we have

$$f(x) \le f(x^{(n)}) + \nabla f(x^{(n)})^{\mathsf{T}}(x - x^{(n)}) + \frac{\rho}{2} ||x - x^{(n)}||^2 \quad (22)$$
$$= g(x|x^{(n)}) \quad (23)$$

 \blacktriangleright x can be updated as

$$x^{(n+1)} = x^{(n)} - (1/\rho)\nabla f(x^{(n)})$$
(24)

$$= x^{(n)} - (1/\rho) (Ax^{(n)} + b)$$
(25)

▶ what if A is positive semidefinite? indefinite? not symmetric?

Jacobi Iterations

• symmetric matrix D and N with positive definite D + N

• we have to find the solution x^* for (D+N)x = -b

note that

$$x^* = \operatorname*{arg\,min}_x f(x) = (1/2) \ x^{\mathsf{T}} (D+N) x + b^{\mathsf{T}} x$$
 (26)

e.g., computing the Newton step for

$$f(x) = \sum_{i=1}^{N} f_i(x_i) + r(Ax - b)$$
(27)

 $x \in {\rm I\!R}^M$, $A \in {\rm I\!R}^{p \times M}$, $b \in {\rm I\!R}^p$, r is a regularization function

• Newton's step Δx_{nt} for f at x is given by

$$\nabla^2 f(x) \Delta x_{\text{nt}} = -\nabla f(x) \tag{28}$$

• the Hessian $\nabla^2 f(x)$ is of the form

$$\nabla^2 f(x) = D + N \tag{29}$$

• identify D and $N \rightarrow$ structured matrix +N

e.g., continues ..

• if N is low rank \rightarrow things can be handled efficiently

▶ if N is not 'sufficiently' low rank \rightarrow apply MM principle

• find $L > ||N||_*$ and majorize $(1/2) x^{\mathsf{T}} N x$

thus f is majorized

more specifically, we have

$$f(x) = (1/2) x^{\mathsf{T}} (D+N)x + b^{\mathsf{T}} x$$
 (30)

$$= (1/2) x^{\mathsf{T}} D x + b^{\mathsf{T}} x + (1/2) x^{\mathsf{T}} N x$$
 (31)

$$\leq (1/2) x^{\mathsf{T}} D x + b^{\mathsf{T}} x + g(x|x^{(n)})$$
 (32)

since $(1/2) \ x^{\mathsf{T}} N x \le g(x|x^{(n)})$, where $g(x|x^{(n)}) = (\rho/2) ||x - x^{(n)}||^2 + x^{(n)\mathsf{T}} N x - (1/2) x^{(n)\mathsf{T}} N x^{\mathsf{T}}$ \blacktriangleright x can be updated as

$$x^{(n+1)} = (D + \rho I)^{-1} \left(\rho x^{(n)} - N x^{(n)} - b \right)$$
(33)

• $(D + \rho I)$ is efficiently invertible \rightarrow has a rich structure

e.g., block diagonal

NONNEGATIVE QUADRATIC PROGRAMMING

The Related Problem

consider the following problem

minimize
$$(1/2)x^{\mathsf{T}}Rx + s^{\mathsf{T}}x$$

subject to $Cx \leq d$

• $x \in \mathbb{R}^N$, positive definite R, suppose it has some structure

$$\blacktriangleright \ C \in {\rm I\!R}^{p \times N} \text{, } d \in {\rm I\!R}^p$$

 \blacktriangleright Lagrangian L is given by

$$L(x,\mu) = (1/2)x^{\mathsf{T}}Rx + s^{\mathsf{T}}x + \mu^{\mathsf{T}}(Cx - d)$$
(34)

 $\mu \in \mathbb{R}^p$

• minimizer $x(\mu)$ of the Lagrangian is given by

$$x(\mu) = -R^{-1}(s + C^{\mathsf{T}}\mu)$$
(35)

 \blacktriangleright dual function h is given by

$$h(\mu) = -(1/2)\mu^{\mathsf{T}}P\mu + q^{\mathsf{T}}\mu + r$$
 (36)

-1

where

$$P = CR^{-1}C^{\mathsf{T}}, \quad q = -(d + CR^{-1}s), \quad r = -\frac{1}{2}s^{\mathsf{T}}R^{-1}s$$

DUAL PROBLEM

dual problem is given by

maximize
$$h(\mu) = -(1/2)\mu^{\mathsf{T}}P\mu + q^{\mathsf{T}}\mu + r$$

subject to $\mu \succeq 0$

how to solve the dual problem?

• interior-point methods ⁴ \rightarrow recall P lacks any structure

 \blacktriangleright not easily implemented for large P

- coordinate descent
- MM principle

 $^{^4 \}text{See} \ \S \ 11.3.1$ of *Convex Optimization* by S. Boyd and L. Vandenberghe, 2004.

COORDINATE DESCENT TO SOLVE THE DUAL

• restrict
$$h$$
 to a line $\mathcal{L}_i = \left\{ \mu^{(n)} + te_i \mid t \in \mathbb{R} \right\}$

• minimize h over the restriction \mathcal{L}_i ⁵

$$\underset{t \in \mathbb{R}}{\text{maximize}} \quad h_i(t) = h(\mu^{(n)} + te_i)$$

 \blacktriangleright we can unfold $h(\mu^{(n)}+te_i)$ in a straightforward manner, i.e., ⁶

$$\begin{aligned} h_i(t) &= -(1/2)(\mu^{(n)} + te_i)^\mathsf{T} P(\mu^{(n)} + te_i) + q^\mathsf{T}(\mu^{(n)} + te_i) \\ &= -(P_{ii}/2) \ t^2 + \left(q_i - \sum_{k=1}^p P_{ik}\mu_i^{(n)}\right) t + \text{irrelevant const.} \end{aligned}$$

⁵Let us first ignore the constrain $\mu \succeq 0$ and assimilate it later.

⁶The constant r is dropped since it is irrelevant.

• compute the derivative h'_i of h to determine t^* , i.e.,

$$-P_{ii}t + \left(q_i - \sum_{k=1}^p P_{ik}\mu_i^{(n)}\right) = 0 \implies t^* = \frac{q_i - \sum_{k=1}^p P_{ik}\mu_i^{(n)}}{p_{ii}}$$

 \blacktriangleright so the $i{\rm th}$ coordinate of current $\mu^{(n)}$ is updated as

$$\mu_i^{(n)} := \mu_i^{(n)} + t^*$$

= $\mu_i^{(n)} + \frac{1}{p_{ii}} \left(q_i - \sum_{k=1}^p P_{ik} \mu_i^{(n)} \right)$

▶ $\mu \succeq 0$ can be assimilated as?

$$\mu_i^{(n)} := \max\left\{0, \mu_i^{(n)} + \frac{1}{p_{ii}} \left(q_i - \sum_{k=1}^p P_{ik} \mu_i^{(n)}\right)\right\}$$

• iterate from i = 1 to i = p and cycles back to i = 1

Algorithm 1 Coordinate Decent

Input: $\mu^{(0)} \succeq 0, n = 0$

1: while a stopping criterion true do

2: **for**
$$i \leftarrow 1$$
 to p **do**
3: $\mu_i^{(n)} \leftarrow \max\left\{0, \mu_i^{(n)} + \frac{1}{p_{ii}}\left(q_i - \sum_{k=1}^p P_{ik}\mu_i^{(n)}\right)\right\}$

4: end for

5:
$$\mu^{(n+1)} = \mu^{(n)}$$
 and $n \leftarrow n+1$

- 6: end while
- 7: return $\mu^{(n+1)}$

• then the solution is given by (35) with $\mu = \mu^{(n+1)}$

MM PRINCIPLE TO SOLVE THE DUAL

recall the objective function ⁷

$$\begin{split} h(\mu) &= -\frac{1}{2} \sum_{i=1}^{N} P_{ii} \mu_{i}^{2} + \sum_{i=1}^{N} q_{i} \mu_{i} \\ &- \frac{1}{2} \sum_{\{i,j|i \neq j, P_{ij} \geq 0\}} P_{ij} \mu_{i} \mu_{j} - \frac{1}{2} \sum_{\{i,j|i \neq j, P_{ij} < 0\}} P_{ij} \mu_{i} \mu_{j} \\ &= -\frac{1}{2} \sum_{i=1}^{N} P_{ii} \mu_{i}^{2} + \sum_{i=1}^{N} q_{i} \mu_{i} \\ &- \frac{1}{2} \sum_{\{i,j|i \neq j, P_{ij} \geq 0\}} P_{ij} \mu_{i} \mu_{j} + \frac{1}{2} \sum_{\{i,j|i \neq j, P_{ij} < 0\}} |P_{ij}| \mu_{i} \mu_{j} \\ &\geq -\frac{1}{2} \sum_{i=1}^{N} P_{ii} \mu_{i}^{2} + \sum_{i=1}^{N} q_{i} \mu_{i} \\ &- \frac{1}{2} \sum_{\{i,j|i \neq j, P_{ij} \geq 0\}} P_{ij} \left[\frac{\mu_{i}^{(n)}}{2\mu_{j}^{(n)}} \mu_{j}^{2} + \frac{\mu_{j}^{(n)}}{2\mu_{i}^{(n)}} \mu_{i}^{2} \right] \\ &+ \frac{1}{2} \sum_{\{i,j|i \neq j, P_{ij} < 0\}} |P_{ij}| \ \mu_{i}^{(n)} \mu_{j}^{(n)} \left[1 + \ln \left(\frac{\mu_{i}}{\mu_{i}^{(n)}} \right) + \ln \left(\frac{\mu_{j}}{\mu_{j}^{(n)}} \right) \right] \\ &= g \left(\mu | \mu^{(n)} \right) \end{split}$$

⁷The constant r is dropped since it is irrelevant.

here the last inequality follows from

$$\mu_i \mu_j \le \frac{\mu_i^{(n)}}{2\mu_j^{(n)}} \ \mu_j^2 + \frac{\mu_j^{(n)}}{2\mu_i^{(n)}} \ \mu_i^2$$

and

$$-\mu_i \mu_j \le -\mu_i^{(n)} \mu_j^{(n)} \left[1 + \ln\left(\frac{\mu_i}{\mu_i^{(n)}}\right) + \ln\left(\frac{\mu_j}{\mu_j^{(n)}}\right) \right]$$

 \blacktriangleright compute the derivative $g'(\ \cdot \ | \mu^{(n)})$ of $g'(\ \cdot \ | \mu^{(n)})$ to yield

$$\begin{split} \left[\sum_{\{i|P_{ki}>0\}} \left(\mu_{i}^{(n)}/\mu_{k}^{(n)}\right)P_{ki}\right]\mu_{k} \\ & -\left[\sum_{\{i|P_{ki}<0\}} \mu_{i}^{(n)}\mu_{k}^{(n)}|P_{ki}|\right] \frac{1}{\mu_{k}} - q_{k} = 0 \\ \Longrightarrow \alpha\mu_{k}^{2} + \beta\mu_{k} + \gamma \quad \text{form} \\ \Longrightarrow \text{ take the positive root as } \mu_{k}^{(n+1)} \end{split}$$

▶ in particular we get

$$\mu_{k}^{(n+1)} = \frac{q_{k} + \sqrt{q_{k}^{2} + 4\left[\sum_{\{i|P_{ki}>0\}} \mu_{i}^{(n)}P_{ki}\right]\left[\sum_{\{i|P_{ki}<0\}} \mu_{i}^{(n)}|P_{ki}|\right]}}{\left[\sum_{\{i|P_{ki}>0\}} \mu_{i}^{(n)}P_{ki}\right]}$$
(37)

iterates can be perform in parallel

Algorithm 2 MM Principle

Input: $\mu^{(0)} \succeq 0, n = 0$

- 1: while a stopping criterion true do
- 2: $\forall k, \mu_k^{(n+1)}$ is computed from (37) and $n \leftarrow n+1$
- 3: end while
- 4: return $\mu^{(n+1)}$
- then the solution is given by (35) with $\mu = \mu^{(n+1)}$
- main differences between MM based algorithm and the coordinate descent?

ARITHMETIC-GEOMETRIC MEAN INEQUALITY

A Majorization to Monomials

weighted arithmetic-geometric mean inequality

$$\prod_{i=1}^{p} x_i^{\alpha_i} \le \sum_{i=1}^{p} \alpha_i x_i \quad \text{for all} \quad x \succeq 0 \tag{38}$$

• α_i are given, $\alpha_i > 0$ ⁸ and $\sum_i \alpha_i = 1$

▶ (38): a majorization to $\prod_{i=1}^{p} x_i^{\alpha_i}$ at $\{\gamma 1 \in \mathbb{R}^p \mid \gamma \in \mathbb{R}_+\}$

• a majorization function to $\prod_{i=1}^p x_i^{\beta_i}$ at

• arbitrary $y \succeq 0$ when $\beta \succ 0$?

⁸If $\alpha_i = 0$, the corresponding x_i is irrelevant.

A GENERAL MAJORIZATION FUNCTION

$$\blacktriangleright$$
 let $\beta_{sum} = \sum_i \beta_i$

▶ substitute $x_i \leftarrow (x_i/y_i)^{\beta_{sum}}$ and $\alpha_i \leftarrow \beta_i/\beta_{sum}$ in (38)

thus, we get

$$\begin{split} \prod_{i=1}^{p} x_{i}^{\beta_{i}} &\leq \left[\prod_{i=1}^{p} y_{i}^{\beta_{i}}\right] \left[\sum_{i=1}^{p} \frac{\beta_{i}}{\beta_{\mathtt{sum}}} \left(\frac{x_{i}}{y_{i}}\right)^{\beta_{\mathtt{sum}}}\right] & \text{for all } x \succeq 0 \\ &= g(x|y) & \text{for all } x \succeq 0 \end{split}$$

A Minorization to Monomials

we rely on the supporting hyperplane inequality

$$\log z \le z - 1 \quad \text{for all} \quad z \in \mathbb{R}_{++} \tag{39}$$

▶ suppose $\beta \succeq 0$ is given, $x_i > 0$

► substitute
$$z = \prod_{i=1}^{p} (x_i/y_i)^{\beta_i}$$
 in (39), i.e.,

$$\prod_{i=1}^{p} x_i^{\beta_i} \ge \prod_{i=1}^{p} y_i^{\beta_i} [1 + \sum_{i=1}^{p} \beta_i \ln (x_i/y_i)] \text{ for all } x \succ 0$$

$$= \prod_{i=1}^{p} y_i^{\beta_i} [1 + \sum_{i=1}^{p} \beta_i \ln x_i - \sum_{i=1}^{p} \beta_i \ln y_i]$$

$$= g(x|y)$$

EXAMPLES

A Majorization to A Signomial

consider the signomial f

$$f(x) = \frac{1}{x_1^3} + \frac{3}{x_1 x_2^2} + x_1 x_2 - \sqrt{x_1 x_2}$$
(40)

majorization function to f at y?

 $1/(x_1x_2^2) \le y_1^2/(3y_2^2x_1^3) + (2y_2)/(3y_1x_2^3)$

• $x_1 x_2 \le (y_2 x_1^2)/(2y_1) + (y_1 x_2^2)/(2y_2)$

 $\sqrt{x_1 x_2} \ge (1/2) \sqrt{y_1 y_2} \left(2 + \ln x_1 + \ln x_2 - \ln y_1 - \ln y_2 \right)$

APPENDICES

Composition with Affine Function

• suppose
$$f : \mathbb{R}^n \to \mathbb{R}$$
 is differentiable

• then define
$$\tilde{f} : \mathbb{R} \to \mathbb{R}$$
 by

$$\tilde{f}(\tau) = f(a + \tau b)$$

is differentiable and $^{\rm 9}$

$$\tilde{f}'(\tau) = \frac{d\tilde{f}(\tau)}{d\tau} = \nabla f(a+\tau b)^{\mathsf{T}}b$$
(41)

 $^{^{9}\}text{see}$ § A.4.2 of Convex Optimization by S. Boyd and L. Vandenberghe, 2004.

Composition with Affine Function

- suppose $f : \mathbb{R}^n \to \mathbb{R}$ is twice differentiable
- then \tilde{f} [cf. (41)] is twice differentiable and ¹⁰

$$\tilde{f}''(\tau) = \frac{d^2 \tilde{f}(\tau)}{d\tau^2} = b^{\mathsf{T}} \nabla^2 f(a + \tau b) b$$
(42)

 $^{^{10}\}text{see}$ § A.4.4 of Convex Optimization by S. Boyd and L. Vandenberghe, 2004.

Newton-Leibniz Formula

▶ recall
$$\tilde{f}(\tau) = f(a + \tau b)$$

let us apply Newton-Leibniz formula ¹¹

$$\tilde{f}(1) = \tilde{f}(0) + \int_0^1 \tilde{f}'(t) dt$$
 (43)

$$f(a+b) = f(a) + \nabla f(a)^{\mathsf{T}}b + \int_0^1 \nabla f(a+tb)^{\mathsf{T}}b \ dt$$
 (44)

¹¹Based on elementary classical analysis.

Taylor with the Integral Remainder

Taylor formula with the integral remainder

• recall
$$\tilde{f}(\tau) = f(a + \tau b)$$

▶ we have ¹²

$$\tilde{f}(1) = \tilde{f}(0) + \tilde{f}'(0) + \int_0^1 \int_0^t \tilde{f}''(\tau) \ d\tau dt$$
(45)

thus from (41) and (42), (45) becomes

$$f(a+b) = f(a) + \nabla f(a)^{\mathsf{T}}b + \int_0^1 \int_0^t b^{\mathsf{T}} \nabla^2 f(a+\tau b)b \, d\tau dt$$
 (46)

¹²Based on elementary classical analysis.